



## **HPC simulation of quantum Computer**



RIKEN R-CCS and supercomputers K 10PFLOPS, 2011-2019



Nobuyasu Ito and Naoki Yoshioka RIKEN Center for Computational Science

> braket: A open-source quantum computer simulator for parallel computer C++ https://github.com/naoki-yoshioka/braket

> > Fugaku 442PFLOPS, 2020-



Fugaku NEXT ~2030?

with quantum-hybrid architecture?



### ISING MACHINE m-TIS

### - first Ising annealer in Japan, second after A. T. Ogielski in 1985 at Bell Lab.



N. Ito, M. Taiji, M. Suzuki, R. Ishibashi, K. Kobayashi, K. Mitsubo and S. Katsura, in "Computer Simulation Studies in Condensed Matter Physics III", edited by D. P. Landau, K. K. Mon and H.-B. Schuettler (Springer--Verlag, 1991) (1991) p.201-203.
M. Taiji, N. Ito and M. Suzuki, Rev. Sci. Intrum. 59 (1988) p.2483

N. Ito, M. Taiji and M. Suzuki, in "Computational Approaches in Condensed Matter Physics", edited by S. Miyashita, M. Imada and H. Takayama (Springer-Verlag, 1992) (1992) p.297-298.





全学講演会の記録

No.4

### R. P. ファインマン

未来の計算様

1985年8月9日講演



Richard P. Feynman

California Institute of Technology Pasadena, California, U.S.A.

It's a great pleasure and an honor to be here as a speaker in memorial for a scientist that I have respected and admired as much as Prof. Nishina. To come to Japan and talk about computers is like giving a sermon to Buddha. But I have been thinking about computers and this is the only subject I could think of when invited to talk.

The first thing I would like to say is what I am not going to talk about. I want to talk about the future computing machines. But the most important possible developments in the future, are things that I will not speak about. For example, there is a great deal of work to try to develop smarter machines, machines which have a better relationship with the humans so that input and output can be made with less effort than the complex programing that's necessary today. This goes under the name often of artificial intelligence, but I don't like that name. Perhaps the unintelligent machines can do even better than the intelligent ones. Another problem is the standardization of programing languages. There are too many languages today, and it would be a good idea to choose just (I hesitate to mention that in Japan, for what will one. happen will be that there will simply be more standard languages; you already have four ways of writing now and

- 1 -

Nishina Memorial Lecture, August 9, 1985, Nishina Foundation and Gakushuin.

学習院大学

### preliminary challenges

### **O**Application of quantum Monte Carlo simulation



#### Hubbard-Stratnovich transformation and multicanonical sampling K. Fischer et al, Intern. J. Modern Phys. C vol. 13 No. 7 (2002) p.917—929. K. Fischer et al, Intern. J. Modern Phys. C vol. 13 No. 7 (2002) p.931—945.

Quantum Circuit from L.M.K. Vandersypen et. al., Nature 414, p.883 (2001): Shor-like circuit with 4 qubits



Fourier transformation

富岳

Fugaku



### **Ostate-vector simulator with GUI: QCAD**







Available online at www.sciencedirect.com

Computer Physics Communications 176 (2007) 121-136

Computer Physics Communications

www.elsevier.com/locate/cpc

#### Massively parallel quantum computer simulator

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### Massively parallel quantum computer simulator, eleven years later



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#### up to 48 gbits

RIKEN Code: https://github.com/naoki-yoshioka/braket

#### up to 36 qbits







Figure B.6: Quantum circuit performing error correction on the top three qubits. The corresponding JUQCS input file is listed in Example: input. Qubits are numbered from zero (top) to four (bottom). Reading from left to right, the first three gates prepare the initial state, the next two (CNOT) gates perform the encoding, the *X* gate on qubit 0 introduces a spin flip error, the next 11 gates detect and correct the error and the last 3 (2) gates on qubit 3 (4) illustrate how to reset a qubit to 0 (1).

#### Program: Quantum assembler – input to simulators

| 1  | QUBITS 5   | 19 | M 3                                     |
|----|--|----|---|
| 2  | ! DEPOLARIZING CHANNEL $p_X = 0.01$ , $p_Y = 0.01$ | 20 | M 4                                     |
| 3  |  | 21 | TOFFOLI 3 4 0                           |
| 4  | H0! initial state                                  | 22 | X 4                                     |
| 5  | Τ 0  | 23 | TOFFOLI 3 4 1                           |
| 6  | H0! initial state                                  | 24 | X 3                                     |
| 7  |  | 25 | X 4                                     |
| 8  | CNOT 0 1 ! encode                                  | 26 | TOFFOLI 3 4 2                           |
| 9  | CNOT 0 2   | 27 | X 3                                     |
| 10 | BEGIN MEASUREMENT                                  | 28 |   |
| 11 |  | 29 | Н 3                                     |
| 12 | !x 1 ! error                                       | 30 | CLEAR 3 ! must be preceeded by Hadamard |
| 13 | $\mathbf{x} = 0 + error$                           | 31 | H 4                                     |
| 14 |  | 32 | SET 4 ! must be preceeded by Hadamard   |
| 15 | CNOT 0 3 / correct                                 | 33 |   |
| 16 | CNOT 1 3   | 34 | BEGIN MEASUREMENT                       |
| 17 | CNOT 0 4   | 35 | GENERATE EVENTS 8192 -1                 |
| 18 | CNOT 2 4   |    |   |





 $| 0 \rangle$  or  $| 1 \rangle$ : orthogonal basis vectors

- $\square$   $n_L n_{L-1} \cdots n_2 n_1$ : a binary representation for n
- State of an L-qubit quantum computer

$$\left|\Phi\right\rangle = \sum_{n=0}^{2^{L}-1} a_{n} \left|n\right\rangle,$$

where  $|n\rangle = |n_L\rangle |n_{L-1}\rangle \cdots |n_2\rangle |n_1\rangle$ ,

$$\sum_{n=0}^{2^{L}-1} |a_{n}|^{2} = 1, \ \langle \Phi | \Phi \rangle = 1$$

 $|a_n|^2$ : probability of obatining n by repeated measurements



### **Requirements for QC simulation**



| F  | ugaku          |  | Memory Iir                          | nits number of q                  | bits.                       |                               |                      |
|--|----------------|--|-------------------------------------|-----------------------------------|-----------------------------|-------------------------------|----------------------|
| Ν  | 2 <sup>N</sup> | double precision $\sim \! 10^{16}$ op. | single precision $\sim \! 10^8$ op. | half precision $\sim$ 10 $^4$ op. | byte precisio $\sim 10^2$ o | n<br>p.                       |                      |
|  |                | (over-spec)                            | (prac                               | tical use)                        | (simple test                | :)                            |                      |
| 10   | 1K             | 16KB                                   | 8KB                                 | 4KB                               | 2KB                         |                               |                      |
| 20   | 1M             | 16MB                                   | 8MB                                 | 4MB                               | 2MB                         |                               |                      |
| 30   | 1G             | 16GB                                   | 8GB                                 | 4GB                               | 2GB                         |                               |                      |
| 36   | 16G            | 1TB                                    | 512GB                               | 256GB                             | 128GB                       | $\leftarrow$ 2007 on Blue Ger | ne/L and Regatta p69 |
| 40   | 1T             | 16TB                                   | 8TB                                 | 4TB                               | 2TB                         |                               |                      |
| 45   | 32T            | 512TB                                  | 256TB                               | 128TB                             | 64TB                        | $\leftarrow$ 2019 on K and Ta | ihu Light            |
| 46   | 64T            | 1PB                                    | 512TB                               | 256TB                             | 128TB                       |                               |                      |
| 47   | 128T           | 2PB                                    | 1PB                                 | 512TB                             | 256TB                       | ← on Fugaku with t            | wo-to-N allocation   |
| 48   | 256T           | 4PB                                    | 2PB                                 | 1PB                               | 512TB                       | ← on Fugaku with p            | acked allocation     |
| 49   | 512T           | 8PB                                    | 4PB                                 | 2PB                               | 1PB                         |                               |                      |
| 50   | 1P             | 16PB                                   | 8PB                                 | 4PB                               | 2PB                         |                               |                      |
| 51   | 2P             | 32PB                                   | 16PB                                | 8PB                               | 4PB                         |                               |                      |
| 52   | 4P             | 64PB                                   | 32PB                                | 16PB                              | 8PB                         |                               |                      |
| $\rightarrow$ Memory transfer rate limits number of gbits. |                | er of qbits.                           | , F                                 | ber node                          | full s                      | ystem                         |                      |
|  |                |  |                                     | in-nod                            | e inter-no                  | ode in-node                   | inter-node           |
|  |                |  |                                     | K: 64GB                           | /s 40GB                     | /s 5.4PB/s                    | 3.4PB/s              |
|  |                |  |                                     | Fugaku: 1TB,                      | /s 40.8G                    | iB/s 155.3PB/s                | 6.2PB/s              |



#### CPU A64FX 7nm FinFET

Architecture Armv8.2-A SVE(512bit SIMD) 48 cores for compute and 4 for OS activities

Normal 2.0GHz (DP:3.072TF, SP: 6.144TF, HP:12.288TF) Boost 2.2GHz(DP:3.3792TF, SP: 6.7584TF, HP:13.5168TF) Cache L1: 64KB 4way >230GB/s(load) >115GB/s(store) L2: 8MB 16way >115GB/s(load) >57GB/s(store) Memory HBM2 8GB X 4=32GiB 1024GB/s

I/O PICe Gen3 x 16 lane

Performance: Stream triad 830GB/s, Dgemm 2.5TF

×8





BoB (Bunch of Blades) 16 Nodes, 8 CMUs

CMU (CPU Memory Unit) 2 Nodes, 2 CPUs

Fugaku: 158,976nodes in 432 rack totally 4.85Peta byte memory with bandwidth 163 PB/s



図 3 A64FX CPU チップの写真



48 Nodes, 3 BoBs

×8



Rack 384 Nodes, 8 Shelves





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**TofuD interconnect** 



Six coordinate axes: X, Y, Z, A, B, C X, Y, Z: the size varies according to the system configuration **A**, **B**, **C**: the size is fixed to  $2 \times 3 \times 2$ ■ Tofu stands for "torus fusion": (X, Y, Z) × (A, B, C)  $X \times Y \times Z \times 2 \times 3 \times 2$ 2

**Max performance** 

each node 40.8 GB/s 2 lanes x 10 ports

full system 6.19 PB hop/s







• 2 MPI processes (M = 1 global qubit)

• process #0 and #1

• Single-qubit gate on the global qubit  $(*|**\cdots)$ 

swap the global qubit and the leftmost local qubit

$$\sigma \mapsto \begin{pmatrix} N & N-1 & N-2 & \cdots & 1 \\ N-1 & N & N-2 & \cdots & 1 \end{pmatrix} \cdot \sigma$$



### How to allocate coefficients to node? $\rightarrow$ qubit swap algorithm

 $\sum$ 

 $n_0, n_1, n_2, n_3, n_4, n_5 = 0, 1$ 

 $a(n_0, n_1, n_2, n_3, n_4, n_5)|n_0n_1n_2n_3n_4n_5>$ 

### node01 a(0,1,0,0,0,0) 0000

| _        |      |                |
|----------|------|----------------|
|          | 0001 | a(0,1,0,0,0,1) |
|          | 0010 | a(0,1,0,0,1,0) |
|          | 0011 | a(0,1,0,0,1,1) |
|          |      |                |
|          | 0100 | a(0,1,0,1,0,0) |
|          | 0101 | a(0,1,0,1,0,1) |
|          | 0110 | a(0,1,0,1,1,0) |
|          | 0111 | a(0,1,0,1,1,1) |
|          |      |                |
|          | 1000 | a(0,1,1,0,0,0) |
|          | 1001 | a(0,1,1,0,0,1) |
| $\sim$   | 1010 | a(0,1,1,0,1,0) |
|          | 1011 | a(0,1,1,0,1,1) |
|          |      |                |
|          | 1100 | a(0,1,1,1,0,0) |
| $\times$ | 1101 | a(0,1,1,1,0,1) |
|          | 1110 | a(0,1,1,1,1,0) |
|          | 1111 | a(0.1.1.1.1.1) |

| 0000 | a(0,0,0,0,0,0) |  |
|------|----------------|--|
| 0001 | a(0,0,0,0,0,1) |  |
| 0010 | a(0,0,0,0,1,0) |  |
| 0011 | a(0,0,0,0,1,1) |  |
|      |                |  |
| 0100 | a(0,0,0,1,0,0) |  |
| 0101 | a(0,0,0,1,0,1) |  |
| 0110 | a(0,0,0,1,1,0) |  |
| 0111 | a(0,0,0,1,1,1) |  |
|      |                |  |
| 1000 | a(0,0,1,0,0,0) |  |
| 1001 | a(0,0,1,0,0,1) |  |
| 1010 | a(0,0,1,0,1,0) |  |
| 1011 | a(0,0,1,0,1,1) |  |
|      | X              |  |
| 1100 | a(0,0,1,1,0,0) |  |
| 1101 | a(0,0,1,1,0,1) |  |
| 1110 | a(0,0,1,1,1,0) |  |
| 1111 | a(0,0,1,1,1,1) |  |

node00

'冨' Fugaku

Case of N=6 qbits on M=4( using 2<sup>4-2</sup>=4 nodes) :  $|\Psi>=$ 

| 0000 | a(1,0,0,0,0,0) | 0000 | a(1,1,0,0,0,0) |
|------|----------------|------|----------------|
| 0001 | a(1,0,0,0,0,1) | 0001 | a(1,1,0,0,0,1) |
| 0010 | a(1,0,0,0,1,0) | 0010 | a(1,1,0,0,1,0) |
| 0011 | a(1,0,0,0,1,1) | 0011 | a(1,1,0,0,1,1) |
|      |                |      |                |
| 0100 | a(1,0,0,1,0,0) | 0100 | a(1,1,0,1,0,0) |
| 0101 | a(1,0,0,1,0,1) | 0101 | a(1,1,0,1,0,1) |
| 0110 | a(1,0,0,1,1,0) | 0110 | a(1,1,0,1,1,0) |
| 0111 | a(1,0,0,1,1,1) | 0111 | a(1,1,0,1,1,1) |
|      |                |      |                |
| 1000 | a(1,0,1,0,0,0) | 1000 | a(1,1,1,0,0,0) |
| 1001 | a(1,0,1,0,0,1) | 1001 | a(1,1,1,0,0,1) |
| 1010 | a(1,0,1,0,1,0) | 1010 | a(1,1,1,0,1,0) |
| 1011 | a(1,0,1,0,1,1) | 1011 | a(1,1,1,0,1,1) |
|      |                |      |                |
| 1100 | a(1,0,1,1,0,0) | 1100 | a(1,1,1,1,0,0) |
| 1101 | a(1,0,1,1,0,1) | 1101 | a(1,1,1,1,0,1) |
| 1110 | a(1,0,1,1,1,0) | 1110 | a(1,1,1,1,1,0) |
| 1111 | a(1,0,1,1,1,1) | 1111 | a(1,1,1,1,1,1) |

Local qubits: allocated to memory-address bits

Global qubits: allocated to node-number bits

R-CCS

 $H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$ 

node10

Apply H 5

















### Swap data in memory blocks ...

| r    | ode00            | r    | node01         |                     | r    | node10           | r    | node11         |
|------|------------------|------|----------------|---------------------|------|------------------|------|----------------|
| 0000 | a(0,0,0,0,0,0)   | 0000 | a(0,1,0,0,0,0) |                     | 0000 | a(1,0,0,0,0,0)   | 0000 | a(1,1,0,0,0,0) |
| 0001 | a(0,0,0,0,0,1)   | 0001 | a(0,1,0,0,0,1) |                     | 0001 | a(1,0,0,0,0,1)   | 0001 | a(1,1,0,0,0,1) |
| 0010 | a(0,0,0,0,1,0)   | 0010 | a(0,1,0,0,1,0) |                     | 0010 | a(1,0,0,0,1,0)   | 0010 | a(1,1,0,0,1,0) |
| 0011 | a(0,0,0,0,1,1)   | 0011 | a(0,1,0,0,1,1) | ~_/                 | 0011 | a(1,0,0,0,1,1)   | 0011 | a(1,1,0,0,1,1) |
| 0100 |                  | 0100 |                | $\sim$              |      |                  | 0100 |                |
| 0100 | a(0,0,0,1,0,0)   | 0100 | a(0,1,0,1,0,0) |                     | 0100 | a(1,0,0,1,0,0)   | 0100 | a(1,1,0,1,0,0) |
| 0101 | a(0,0,0,1,0,1)   | 0101 | a(0,1,0,1,0,1) |                     | 0101 | a(1,0,0,1,0,1)   | 0101 | a(1,1,0,1,0,1) |
| 0110 | a(0,0,0,1,1,0)   | 0110 | a(0,1,0,1,1,0) |                     | 0110 | a(1,0,0,1,1,0)   | 0110 | a(1,1,0,1,1,0) |
| 0111 | a(0,0,0,1,1,1)   | 0111 | a(0,1,0,1,1,1) | $\times$            | 0111 | a(1,0,0,1,1,1)   | 0111 | a(1,1,0,1,1,1) |
|      |                  |      |                |                     |      |                  |      |                |
| 1000 | a(0,0,1,0,0,0)   | 1000 | a(0,1,1,0,0,0) |                     | 1000 | a(1,0,1,0,0,0)   | 1000 | a(1,1,1,0,0,0) |
| 1001 | a(0,0,1,0,0,1)   | 1001 | a(0,1,1,0,0,1) |                     | 1001 | a(1,0,1,0,0,1)   | 1001 | a(1,1,1,0,0,1) |
| 1010 | a(0,0,1,0,1,0)   | 1010 | a(0,1,1,0,1,0) | $\sim$              | 1010 | a(1,0,1,0,1,0)   | 1010 | a(1,1,1,0,1,0) |
| 1011 | a(0,0,1,0,1,1)   | 1011 | a(0,1,1,0,1,1) | $\langle \ \rangle$ | 1011 | a(1,0,1,0,1,1)   | 1011 | a(1,1,1,0,1,1) |
|      |                  |      |                |                     |      |                  |      |                |
| 1100 | a(0,0,1,1,0,0)   | 1100 | a(0,1,1,1,0,0) |                     | 1100 | a(1,0,1,1,0,0)   | 1100 | a(1,1,1,1,0,0) |
| 1101 | a(0,0,1,1,0,1)   | 1101 | a(0,1,1,1,0,1) | < /                 | 1101 | a(1,0,1,1,0,1)   | 1101 | a(1,1,1,1,0,1) |
| 1110 | a(0,0,1,1,1,0) 🕴 | 1110 | a(0,1,1,1,1,0) | $\sim$              | 1110 | a(1,0,1,1,1,0) 🕴 | 1110 | a(1,1,1,1,1,0) |
| 1111 | a(0,0,1,1,1,1)   | 1111 | a(0,1,1,1,1,1) |                     | 1111 | a(1,0,1,1,1,1)   | 1111 | a(1,1,1,1,1,1) |
|      | $\sim$           |      |                |                     |      |                  |      |                |





node *n<sub>0</sub>n<sub>1</sub>n<sub>2</sub>...* 

| a(0,0,0,*) | a(0,0,1,*) |
|------------|------------|
| a(0,1,0,*) | a(0,1,1,*) |
| a(1,0,0,*) | a(1,0,1,*) |
| a(1,1,0,*) | a(1,1,1,*) |

1 bit swap: swap of 1/2

2 bit swap: swap of 3/4

3 bit swap: swap of 7/8

...



Table 1: Overview of the computer systems used for benchmarking. The IBM Blue Gene/Q JUQUEEN [3] (decomissioned), JURECA [4] and JUWELS are located at the Jülich Supercomputing Centre in Germany, the K computer of the RIKEN Center for Computational Science in Kobe, Japan, and the Sunway TaihuLight [5] at the National Supercomputer Center in Wuxi, China. The row "# qubits" gives the maximum number of qubits N that can be simulated with JUQCS-A (JUQCS-E). At the time of running the benchmarks on JUWELS, the maximum number of qubits N was limited to 43 (40).

|                      | JUQUEEN                  | K computer                   | Sunway TaihuLight               | JURECA-CLUSTER           | JUWELS                           |
|----------------------|--------------------------|------------------------------|---------------------------------|--------------------------|----------------------------------|
| CPU                  | IBM PowerPC<br>A2        | eight-core SPARC64<br>VIIIfx | SW26010 manycore<br>64-bit RISC | Intel Xeon<br>E5-2680 v3 | Dual Intel Xeon<br>Platinum 8168 |
| clock frequency      | 1.6 GHz                  | 2.0 Ghz                      | 1.45 GHz                        | 2.5 GHz                  | 2.7 GHz                          |
| memory/node          | 16 GB                    | 16 GB                        | 32 GB                           | 128 GB                   | 96 GB                            |
| # threads/core used  | 1 - 2                    | 8                            | 1                               | 1 - 2                    | 1 - 2                            |
| # cores used         | 1 - 262144               | 2 - 65536                    | 1 - 131072                      | 1 - 6144                 | 1 - 98304                        |
| # nodes used         | 1 – 16384                | 2 - 65536                    | 1 - 32768                       | 1 - 256                  | 1 - 2048                         |
| # MPI processes used | 1 - 524288               | 2 - 65536                    | 1 - 131072                      | 1 - 1024                 | 1 - 2048                         |
| # qubits             | 46 (43)                  | 48 (45)                      | 48 (45)                         | 43 (40)                  | 46 (43)                          |
| X                    | No.1 in 2011<br>10PFLOPS | No.5 in 2013<br>5PFLOPS      | No.1 in 2016                    | X                        | X                                |

#### **Double precision**



Normalied elapsed time by elapsed time of 32 gbits case

#### Normalization factors:

- JUQUEEN 1.2sec
- K 1.0sec
- Sunway TaihuLight 7.7sec
- JURECA 1.9sec • JUWELS 1.3sec

H. De Raedt, F. Jin, D. Wilsch, M. Nocon, N. Yoshioka, N. Ito, S. Yuan and K. Michielsen, Comp. Phys. Comm. Vol.237 (2019) p.47-61







### **Performance: one Hadamard operation for all qbits**

### more qbits with reduced precision using one-byte coding



Normalied elapsed time by elapsed time of 35 qbits case

#### Normalization factors:

| JUQUEEN       | 2.7sec       |
|---------------|--------------|
| K             | 3.8sec       |
| Supway Taibul | ight 10 Ococ |

- Sunway TaihuLight 19.9sec
- JURECA 2.4sec
- JUWELS 2.2sec

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- 1 unit (M = 0 global qubit), K = 3 unit qubits,  $n_{\rm u} = 3$  MPI processes
  - **process** #0: u = 000, 001, 010, #1: u = 011, 100, 101, #2: u = 110, 111
  - apply single-qubit gate on rightmost unit qubit (\* \* \* |···)
    - swap the unit qubit and the leftmost local qubit





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### Performance of Hadamard benchmark(once) on the "Fugaku"





cf on K, 35 qubits: 1.2 sec 40 qubits: 2.0 sec

- no unit : basic parallelization using 2<sup>M</sup> processes
- others : non 2<sup>M</sup> parallelization





### Limitations of the qubit-swap implementations

### Number of coefficients on each node is $2^n$ .

|        | memory per node                   | available memory          |  |
|--------|-----------------------------------|---------------------------|--|
| Κ      | 16GB                              | 8 GB                      |  |
| Fugaku | 32GB                              | 16 GB                     |  |
| Number | of used nodes is 2 <sup>m</sup> . |                           |  |
|        | number of nodes                   | available nodes           |  |
| Κ      | 88,128                            | 65,536 = 2 <sup>16</sup>  |  |
| Fugaku | 158,976                           | 131,072 = 2 <sup>17</sup> |  |

### Remove these limitations by allocating more qubits

Pack more coefficients to each node:

Local set: use smaller number of local qubits 0.5GB allocate many local sets to each node 53 local sets using 26.5GB

- → 4PB using 158,276 nodes of 158,976 nodes of Fugaku
   48 qubits(double), 49(single), 50(half) and 51(byte)
- Or "unit" qubits between local and global qubits Coefficients of (unit+local) qubits are in some nodes or processes.

Limitations of memory band-width

In node: 1TB/s(Fugaku) ~100GB(x86 Workstation) Internode: 40.8GB/s(Fugaku)

one load and one store per operatio 16GB(30 qubits in double)/node  $\rightarrow$  16GBx2/1TB = 31 msec one bit-swap per operation

8GB transfer  $\rightarrow$  8TBx2/40.8GB = 392 msec

sub-second per operation  $\rightarrow \sim 10^5$  operations in 10 hours

Remove these limitations by optimal allocation to increase circuit depth

0.5GB(26 qubits in single)/node → 0.5GBx2/1TB = 1 msec internode 0.25GBx2/40.8GB = 12 msec

10 ms per operation  $\rightarrow \sim 10^6$  operations in 10 hours







various number of qbit and number of local bit The cost unit is sending and receiving one page, i.e. 2^(I-2) amplitudes, for each nodes.

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### **Summary**



- Algorithm and program for massively parallel supercomputers are showns.
- On Fugaku, state-vector simulation can reach
   48 qubits in double precision and 51 qubits in byte precision.

### Perspectives

- Further tuning, mainly for data transfer, will improve performance.
- Applications to development of quantum computer and to test quantum algorithms have been started jointly with RIKEN quantum computer center(RQC).
- Optimization of execution speed.

braket: A open-source quantum computer simulator for parallel computer C++, two-to-N implementation, no unit qubit https://github.com/naoki-yoshioka/braket





# Software environment for classical-quantum hybrid computation





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# Supercomputer Fugaku at RIKEN R-CCS in Kobe (© RIKEN) CCP2023 - 34th IUPAP Conference on Computational Physics

- Date: August 4 (Fri) 8 (Tue), 2023
- Venue: Kobe International Conference Center, Kobe Port Island, Kobe, Japan

