Challenges for Scaling: Co-Design for Memory Bottleneck, Power and Miniaturization

Group B

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Synthesis

- Large amount of data, that is mostly irregular and at times need to be processed at the edge, poses new challenges for scaling:
- Need for programming, architecture and power improvements.
 - Memory Bottlenecks
 - Portability (Miniaturization and Power efficiency)
 - Programmer productivity

Motivations

- Democratizing Compute: (Bioinformatics & Smart Medical Systems)
 - Dataflow in Scientific Workflows
 - Intelligent Medical Systems Real Time Processing
- Scientific Simulations: (Quantum physics & Weather Forecasting)
 - Multi Precision Arithmetics
 - Data Assimilation & Learning
- Memory Acceleration: (Graph Processing & Machine Intelligence)
 - Non-von Neumann Architectures
 - Continuum Computer Architecture
 - Neuromorphic

Problem Domain: Scientific Workflow with Containers

Omics (genomics, metabolomics, proteomics), machine learning pipelines, virtual drug screening



Scientific workflows

Problem: network contention



Solution: Dataflow programming model



Memory is used for intermediate results.

How move data to/from containers?

- UNIX pipes
- Memory-mapped files
- Tmpfs

High-level API hides parallel computing challenges

• User productivity

Scales on cloud and commodity HW

Problem Domain: Biomedical Diagnosis

- Processing massive streams of data is an important problem in Biomedical diagnosis systems.
 - Biomedical diagnosis involves real time signal processing
 - A large number of transducers used, which generate massive data
 - Signal processing algorithms require huge memory to store pre computed coefficients
 - Accessing memory makes system performance slow : a bottleneck in real-time diagnosis

Example -

3D Ultrasound imaging requires 50 GB LUT (Lookup tables) space

Solution: Biomedical Diagnosis



- Exploring sparsity of the data : compressive sensing
- Customized hardware : parallel computing
- On the fly computation : reduced memory access

Problem Domain: Quantum Physics

Numerical calculation for quantum physics

1 What is the presence problem about quantum physics ?

2 Making program for numerical calculation

Considering computation time and capacity of files

Einstein equation

$$R_{\mu\nu} - \frac{1}{2}g_{\mu\nu}R = \frac{8\pi G}{c^4}T_{\mu\nu}$$

Schrodinger equation





Problem Domain: Weather Forecasting

Data size issues in data assimilation

Observational data size issues:

Real-time finescale weather forecast requires much observational data input

- conventional techniques (radar, satellites) with higher resolution
- new data sources (vehicles, portable devices)

Fast computation and data transfer are both essential

Possible solutions:

- improved pre-processing schemes



Problem Domain: Linear Algebra

Multi precision arithmetic

Double-Double and Quad-Double arithmetic uses the combinations of double precision numbers. # of operations would become large.

In the conventional laptop computer,

- Without parallelization, a kernel (BLAS 1 2 3) is computation bottleneck.
- With parallelization(FMA, SIMD, OpenMP), some kernels are memory bottleneck.

Parallelization have memory performance constraint for some multi precision kernels.

Problem Domain: Machine Learning

Memory Access - Bottleneck for DL applications.



- 1. DRAM access: Data movement DRAM to ALU is expensive.
- 2. Mapping data-flow over the architecture: Memory hierarchy to computation units.
- 3. For DL application training and inferencing, loading huge data for training affects the training time, which may be critical for many real-time applications.

Solution: Machine Learning

- 1. Data compression to reduce the storage and movement.
- 2. Network pruning e.g based on magnitude of weights.
- 3. Reduce precision for computation: (Floating point -> Fixed point): 8 bit int used in (Google TPU).
 - a. Binary weight, ternary weight..
 - b. Non linear quantization (Log-domain)
- 4. Improve the reuse of data and local (computational) accumulation.
- 5. Exploit sparsity in the computation map: skip memory access and compute for zero.
- 6. Reduce operation while mapping DNN to matrix multiplication, example using FFT.
- 7. On-chip memory partition, putting memory and processor on same silicon substrate, increase the memory Bandwidth.
- 8. Moving from temporal architecture (SIMD) (MEM-> REG File -> ALU -> control) to Spatial architecture (more advanced for memory accessing) (MEM -> ALU).
- 9. Advance memory techniques: Stacked DRAMs and non-volatile memories.
- 10. Explore possibility of neuromorphic computing with asynchronous operation.

Problem Domain: Graph Processing

- Graph processing generally involves:
 - Low FLOP to Byte ratio
 - Irregular data access pattern
- Bulk Synchronous Model (BSP) leads to under exploitation of the large inherent parallelism that is naturally available in graph structures.
- Think like a Vertex, asynchronously:
- Send an active message asynchronous (fire-and-forget).
- No DAG. Because there could be cycles in the graph.
- We implement Dijkstra–Scholten algorithm for termination detection



Problem Domain: Graph Processing



Listing 1: Asynchronous SSSP



Presents both behaviors of Strong and Weak Scaling: Transcendental Scaling



Problem Domain: Graph Processing

- Continuum Compute Architecture is a new class of non von Neumann architectures.
- Offers fine grain parallelism.
- Small compute cells organized such that it creates an active memory.
- Low Power
- Less space footprint



Conclusion

- New Challenges posed by Big Data
 - Irregular memory access
 - Memory bottleneck
 - Latency sensitive
 - Low Power requirements
- Solutions:
 - 3D stacked Memory
 - Non-von Neumann architectures: send work/compute to memory and process there
 - Custom hardware for inference (and other compute) -> less power and less areas footprint, critical for portability
 - Dataflow-oriented workflows
 - Programmer productivity
 - Auto optimizations (lazy evaluation, concurrency, locality optimization)